

AMS526: Numerical Analysis I (Numerical Linear Algebra)

Lecture 14: Gaussian Elimination and LU Factorization

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Gaussian Elimination and LU Factorization

- Gaussian elimination was covered intensively in AMS 505
- Gaussian elimination can be viewed as “triangular triangularization” of nonsingular $\mathbf{A} \in \mathbb{C}^{m \times m}$

$$\underbrace{L_{m-1} \cdots L_2 L_1}_{L^{-1}} \mathbf{A} = \mathbf{U}$$

analogous to Householder QR factorization of matrix $\mathbf{A} \in \mathbb{C}^{m \times n}$

$$\underbrace{Q_n \cdots Q_2 Q_1}_{Q^*} \mathbf{A} = \mathbf{R}$$

- Example of LU factorization of 4×4 matrix \mathbf{A}

$$\begin{array}{c}
 \xrightarrow{L_1} \\
 \underbrace{\begin{bmatrix} \times & \times & \times & \times \\ 0 & \times & \times & \times \\ 0 & \times & \times & \times \\ 0 & \times & \times & \times \end{bmatrix}}_{L_1 \mathbf{A}}
 \end{array}
 \xrightarrow{L_2}
 \begin{array}{c}
 \underbrace{\begin{bmatrix} \times & \times & \times & \times \\ & \times & \times & \times \\ & 0 & \times & \times \\ & 0 & \times & \times \end{bmatrix}}_{L_2 L_1 \mathbf{A}}
 \end{array}
 \xrightarrow{L_3}
 \begin{array}{c}
 \underbrace{\begin{bmatrix} \times & \times & \times & \times \\ & \times & \times & \times \\ & & \times & \times \\ & & 0 & \times \end{bmatrix}}_{L_3 L_2 L_1 \mathbf{A}}
 \end{array}$$

Forming L

- Luckily, the L matrix contains the *multipliers* $l_{jk} = x_{jk}/x_{kk}$

$$L = L_1^{-1} L_2^{-1} \cdots L_{m-1}^{-1} = \begin{bmatrix} 1 & & & & & \\ l_{21} & 1 & & & & \\ l_{31} & l_{32} & 1 & & & \\ \vdots & \vdots & \ddots & \ddots & & \\ l_{m1} & l_{m2} & \cdots & l_{m,m-1} & 1 & \end{bmatrix}$$

and is said to be a *unit lower triangular matrix*

- First, $L_k^{-1} = I + l_k e_k^*$, because $e_k^* l_k = 0$ and $(I - l_k e_k^*)(I + l_k e_k^*) = I - l_k e_k^* l_k e_k^* = I$
- Second, $L_1^{-1} L_2^{-1} \cdots L_{k+1}^{-1} = I + \sum_{j=1}^{k+1} l_j e_j^*$, since (prove by induction) $(I + \sum_{j=1}^k l_j e_j^*)(I + l_{k+1} e_{k+1}^*) = I + \sum_{j=1}^{k+1} l_j e_j^* + \sum_{j=1}^k l_j (e_j^* l_{k+1}) e_{k+1}^*$ where $e_j^* l_{k+1} = 0$ for $j < k + 1$
- In other words, L is “union” of $L_1^{-1}, L_2^{-1}, \dots, L_{m-1}^{-1}$

Gaussian Elimination without Pivoting

- Factorize $\mathbf{A} \in \mathbb{C}^{m \times m}$ into $\mathbf{A} = \mathbf{L}\mathbf{U}$

Gaussian elimination without pivoting

$$\mathbf{U} = \mathbf{A}, \mathbf{L} = \mathbf{I}$$

for $k = 1$ to $m - 1$

for $j = k + 1$ to m

$$\ell_{jk} = u_{jk}/u_{kk}$$

$$u_{j,k:m} = u_{j,k:m} - \ell_{jk}u_{k,k:m}$$

- Flop count $\sim \sum_{k=1}^m 2(m-k)(m-k) \sim 2 \sum_{k=1}^m k^2 \sim 2m^3/3$
- In actually, \mathbf{L} often overwrites lower-triangular part of \mathbf{A} and \mathbf{U} overwrites upper-triangular part of \mathbf{A}
- Question: What if u_{kk} is 0? Answer: The algorithm would break.

Partial Pivoting

- At step k , we divide by u_{kk} , which would break if u_{kk} is 0 (or close to 0), which can happen even if \mathbf{A} is nonsingular

$$\begin{bmatrix} \times & \times & \times & \times & \times \\ & x_{kk} & \mathbf{x} & \mathbf{x} & \mathbf{x} \\ & \times & \times & \times & \times \\ & \times & \times & \times & \times \\ & \times & \times & \times & \times \end{bmatrix} \rightarrow \begin{bmatrix} \times & \times & \times & \times & \times \\ & x_{kk} & \times & \times & \times \\ & 0 & \mathbf{x} & \mathbf{x} & \mathbf{x} \\ & 0 & \mathbf{x} & \mathbf{x} & \mathbf{x} \\ & 0 & \mathbf{x} & \mathbf{x} & \mathbf{x} \end{bmatrix}$$

- However, any nonzero entry in k th column below diagonal can also be used as *pivot*

$$\begin{bmatrix} \times & \times & \times & \times & \times \\ & \times & \times & \times & \times \\ & \times & \times & \times & \times \\ & x_{ik} & \mathbf{x} & \mathbf{x} & \mathbf{x} \\ & \times & \times & \times & \times \end{bmatrix} \rightarrow \begin{bmatrix} \times & \times & \times & \times & \times \\ & 0 & \mathbf{x} & \mathbf{x} & \mathbf{x} \\ & 0 & \mathbf{x} & \mathbf{x} & \mathbf{x} \\ & x_{ik} & \times & \times & \times \\ & 0 & \mathbf{x} & \mathbf{x} & \mathbf{x} \end{bmatrix}$$

and we permute (interchange) row i with row k

- In general, we take nonzero entry with largest absolute value

More on Partial Pivoting

- k th step of Gaussian elimination of partial pivoting

$$\begin{array}{ccc}
 \begin{bmatrix} \times & \times & \times & \times \\ & \times & \times & \times \\ & & x_{ik} & \mathbf{x} & \mathbf{x} \\ & & \times & \times & \times \end{bmatrix} & \xrightarrow{P_k} & \begin{bmatrix} \times & \times & \times & \times \\ & x_{kk} & \mathbf{x} & \mathbf{x} \\ & \mathbf{x} & \mathbf{x} & \mathbf{x} \\ & \times & \times & \times \end{bmatrix} & \xrightarrow{L_k} & \begin{bmatrix} \times & \times & \times & \times \\ & x_{kk} & \times & \times \\ & 0 & \mathbf{x} & \mathbf{x} \\ & 0 & \mathbf{x} & \mathbf{x} \end{bmatrix} \\
 \text{Pivot selection} & & \text{Row interchange} & & \text{Elimination}
 \end{array}$$

and we interchange row i with row k

- In terms of matrices, it becomes $\underbrace{L_{m-1}P_{m-1} \cdots L_2P_2L_1P_1}_{L^{-1}P} A = U$

- $P = P_{m-1} \cdots P_2P_1$ and $L = (L'_{m-1} \cdots L'_2L'_1)^{-1}$, where $L'_k = P_{m-1} \cdots P_{k+1}L_kP_{k+1}^{-1} \cdots P_{m-1}^{-1}$

- It is easy to verify that

$$L_{m-1}P_{m-1} \cdots L_2P_2L_1P_1 = (L'_{m-1} \cdots L'_2L'_1) (P_{m-1} \cdots P_2P_1)$$

- $L'_k = I - P_{m-1} \cdots P_{k+1}\ell_k e_k^*$ and L is “union” of $(L'_k)^{-1} \equiv I + P_{m-1} \cdots P_{k+1}\ell_k e_k^*$

Algorithm of Gaussian Elimination with Partial Pivoting

- Factorize $A \in \mathbb{C}^{m \times m}$ into $PA = LU$

Gaussian elimination with partial pivoting

$$U = A, L = I, P = I$$

for $k = 1$ to $m - 1$

$$i \leftarrow \arg \max_{i \geq k} |u_{ik}|$$

$$u_{k,k:m} \leftrightarrow u_{i,k:m}$$

$$l_{k,1:k-1} \leftrightarrow l_{i,1:k-1}$$

$$p_{k,:} \leftrightarrow p_{i,:}$$

for $j = k + 1$ to m

$$l_{jk} = u_{jk}/u_{kk}$$

$$u_{j,k:m} = u_{j,k:m} - l_{jk}u_{k,k:m}$$

- Question: What if u_{kk} is 0?
- Flop count $\sim \sum_{k=1}^m 2(m-k)(m-k) \sim 2 \sum_{k=1}^m k^2 \sim 2m^3/3$, same as without pivoting

An Alternative Implementation

- In practice, L and U overwrite A and P is represented by a vector

Gaussian elimination with partial pivoting (alternative)

$$p = [1, 2, \dots, m];$$

for $k = 1$ to $m - 1$

$$i \leftarrow \arg \max_{i \geq k} |a_{ik}|$$

$$a_{k,1:m} \leftrightarrow a_{i,1:m}$$

$$p_k \leftrightarrow p_i$$

$$a_{k+1:m,k} \leftarrow a_{k+1:m,k} / a_{k,k}$$

$$A_{k+1:m,k+1:m} \leftarrow A_{k+1:m,k+1:m} - a_{k+1:m,k} * a_{k,k+1:m}$$

- Using LU factorization to solve $Ax = b$:
 - 1 $PA = LU$; (LU factorization with partial pivoting)
 - 2 $Ly = P^T b$; (Forward substitution)
 - 3 $Ux = y$; (Back substitution)

Complete Pivoting

- More generally, we can use any nonzero entry
- In theory, any nonzero entry (i, j) , $i \geq k, j \geq k$

$$\begin{bmatrix} \times & \times & \times & \times \\ & \times & \times & \times \\ & \mathbf{x} & x_{ij} & \mathbf{x} \\ & \times & \times & \times \end{bmatrix} \rightarrow \begin{bmatrix} \times & \times & \times & \times \\ & \mathbf{x} & 0 & \mathbf{x} \\ & \times & x_{ij} & \times \\ & \mathbf{x} & 0 & \mathbf{x} \end{bmatrix}$$

and we then permute row i with row k , column j with column k

- In matrix operations, it can be expressed as

$$\underbrace{L_{m-1}P_{m-1} \cdots L_2P_2L_1P_1}_{L^{-1}P} \underbrace{AQ_1Q_2 \cdots Q_{m-1}}_Q = U$$

- Therefore, $PAQ = LU$ where $P = P_{m-1} \cdots P_2P_1$ and $L = (L'_{m-1} \cdots L'_2L'_1)^{-1}$
- However, complete pivoting is typically not used in practice because it increases cost in search of pivot and complexity of implementation