

# AMS526: Numerical Analysis I (Numerical Linear Algebra)

## Lecture 25: Preconditioning; Other Krylov Subspace Methods

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# Outline

1 Preconditioning

2 GMRES and Other Krylov Subspace Methods

# Preconditioning

- Motivation: Convergence of iterative methods heavily depends on the condition number of the equation
- Main idea of preconditioning is to introduce a nonsingular matrix  $M$  such that  $M^{-1}A$  has a smaller condition number. Thereafter, solve

$$M^{-1}Ax = M^{-1}b,$$

which has the same solution as  $Ax = b$

- Criteria of  $M$ 
  - ▶ Good approximation of  $A$
  - ▶ Ease of inversion
- However, if  $A$  is symmetric,  $M^{-1}A$  breaks symmetry. How to resolve this problem?

## Preconditioning of CG

- When preconditioning a symmetric matrix, use SPD matrix  $M$
- Let its Cholesky factorization be  $M = RR^*$ . Solve

$$R^{-1}AR^{-T}(R^T x) = R^{-1}b$$

- In practice, algorithm can be organized so that only  $M^{-1}$  (instead of  $R^{-1}$ ) appears

Algorithm: Preconditioned Conjugate Gradient Method

$$x_0 = 0, r_0 = b, p_0 = M^{-1}r_0, z_0 = p_0$$

for  $n = 1$  to  $1, 2, 3, \dots$

$$\alpha_n = (r_{n-1}^T z_{n-1}) / (p_{n-1}^T A p_{n-1})$$

step length

$$x_n = x_{n-1} + \alpha_n p_{n-1}$$

approximate solution

$$r_n = r_{n-1} - \alpha_n A p_{n-1}$$

residual

$$z_n = M^{-1}r_n$$

preconditioning

$$\beta_n = (r_n^T z_n) / (r_{n-1}^T z_{n-1})$$

improvement this step

$$p_n = r_n + \beta_n p_{n-1}$$

search direction

## Commonly Used Preconditioners

- **Jacobi preconditioning:**  $M = \text{diag}(\mathbf{A})$ . Very simple and cheap, might improve certain problems but usually insufficient
- **Block-Jacobi preconditioning:** Let  $M$  be composed of block-diagonal instead of diagonal.
- **Classical iterative methods:** Precondition by applying one step of Jacobi, Gauss-Seidel, SOR, or SSOR
- **Incomplete factorizations:** Perform Gaussian elimination but ignore fill
- **Multigrid (coarse-grid approximations):** For a PDE discretized on a grid, a preconditioner can be formed by transferring the solution to a coarser grid, solving a smaller problem, then transferring back. This is the most efficient approach if applicable

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# Minimizing Residual

- CG only works for SPD matrices.
- There have been many proposed extensions to nonsymmetric matrices, GMRES, BiCG, etc.
- GMRES (Generalized Minimal RESiduals) is one of most well known
- The basic idea is the following
  - ▶ Find  $\mathbf{x}_n \in \mathcal{K}_n$  that minimizes  $\|\mathbf{r}_n\| = \|\mathbf{b} - \mathbf{A}\mathbf{x}_n\|$
  - ▶ This can be viewed as a least squares problem: Find a vector  $\mathbf{c}$  s.t.  $\|\mathbf{A}\mathbf{K}_n\mathbf{c} - \mathbf{b}\|$  is minimized, where  $\mathbf{K}_n$  is the  $m \times n$  Krylov matrix composed of basis vectors of  $\mathcal{K}_n$
  - ▶ Orthogonal basis is often used, produced by Arnoldi iteration

## Review: Arnoldi Iteration

- Let  $Q_n = [q_1 | q_2 | \cdots | q_n]$  be  $m \times n$  matrix with first  $n$  columns of  $Q$  and  $\tilde{H}_n$  be  $(n+1) \times n$  upper-left section of  $H$
- Start by picking a random  $q_1$  and then determine  $q_2$  and  $\tilde{H}_1$
- The  $n$ th columns of  $AQ_n = Q_{n+1}\tilde{H}_n$  can be written as

$$Aq_n = h_{1n}q_1 + \cdots + h_{nn}q_n + h_{n+1,n}q_{n+1}$$

Algorithm: Arnoldi Iteration

given random nonzero  $b$ , let  $q_1 = b/\|b\|$

for  $n = 1$  to  $1, 2, 3, \dots$

$$v = Aq_n$$

for  $j = 1$  to  $n$

$$h_{jn} = q_j^* v$$

$$v = v - h_{jn}q_j$$

$$h_{n+1,n} = \|v\|$$

$$q_{n+1} = v/h_{n+1,n}$$

## Minimal Residual with Orthogonal Basis

- Let  $Q_n$  be Krylov matrix whose columns  $q_1, q_2, \dots$  span the successive Krylov subspaces
- Instead of find  $x_n = \mathcal{K}_n c$ , find  $x_n = Q_n y$  which minimizes  $\|AQ_n y - b\|$
- For Arnoldi iteration, we showed that  $AQ_n = Q_{n+1} \tilde{H}_n$ , so

$$\|Q_{n+1} \tilde{H}_n y - b\| = \text{minimum}$$

- Left multiplication by  $Q_{n+1}^*$  does not change the norm, so

$$\|\tilde{H}_n y - Q_{n+1}^* b\| = \text{minimum}$$

- Finally, by construction,  $Q_{n+1}^* b = \|b\| e_1$ , so

$$\|\tilde{H}_n y - \|b\| e_1\| = \text{minimum.}$$

# The GMRES Algorithm

Algorithm: GMRES

$$q_1 = b/\|b\|$$

for  $n = 1$  to  $1, 2, 3, \dots$

    Step  $n$  of Arnoldi iteration

    Find  $y$  to minimize  $\|\tilde{H}_n y - \|b\|e_1\| = \|r_n\|$

$$x_n = Q_n y$$

- The residual  $\|r_n\|$  does not need to be computed explicitly from  $x_n$
- Least squares problem has Hessenberg structure, solve with QR factorization of  $\tilde{H}_n$  (computed by updating the factorization of  $\tilde{H}_{n-1}$ )
- Memory and cost grow with  $n$ . Restart the algorithm by clearing accumulated data (might stagnate the method)

# Convergence of GMRES

- GMRES converges monotonically and it converges after at most  $m$  steps
- The residual  $\mathbf{r}_n = p_n(\mathbf{A})\mathbf{b}$ , where  $p_n \in P_n$  is a degree  $n$  polynomial with  $p(0) = 1$ , so GMRES also finds a minimizing polynomial such that  $\|p_n(\mathbf{A})\mathbf{b}\| = \text{minimum}$
- Based on a polynomial analysis, diagonalizable  $\mathbf{A} = \mathbf{V}\mathbf{\Lambda}\mathbf{V}^{-1}$  converges as

$$\frac{\|\mathbf{r}_n\|}{\|\mathbf{b}\|} \leq \kappa(\mathbf{V}) \inf_{p_n \in P_n} \sup_{\lambda_i \in \Lambda(\mathbf{A})} |p_n(\lambda_i)|$$

- In other words, if  $\mathbf{A}$  is not far from normal (i.e., eigenvectors are nearly orthogonal), and if properly normalized degree  $n$  polynomials can be found whose size on the spectrum  $\Lambda(\mathbf{A})$  decreases quickly with  $n$ , then GMRES converges quickly

## Other Krylov Subspace Methods

- CG on the Normal Equations (CGN)
  - ▶ Solve  $\mathbf{A}^T \mathbf{A} \mathbf{x} = \mathbf{A}^T \mathbf{b}$  using Conjugate Gradients
  - ▶ Poor convergence due to squared condition number (i.e.,  $\kappa(\mathbf{A}^T \mathbf{A}) = \kappa(\mathbf{A})^2$ ).
- BiConjugate Gradients (BiCG)
  - ▶ Makes residuals orthogonal to another Krylov subspace, based on  $\mathbf{A}^*$
  - ▶ Memory requirements only constant number of vectors
  - ▶ Convergence sometimes comparable to GMRES, but unpredictable
- Conjugate Gradients Squared (CGS)
  - ▶ Avoids multiplication by  $\mathbf{A}$ , sometimes twice as fast convergence
- Quasi-Minimal Residuals (QMR) and Stabilized BiCG (Bi-CGSTAB)
  - ▶ Variants of BiCG with more regular convergence